A framework for a modular multi-concept lexicographic closure semantics (an abriged report)¹

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Aim of the talk

We define a *modular multi-concept extension of the lexicographic closure semantics* for defeasible description logics with typicality.

The idea is that of *distributing the defeasible properties of concepts into different* modules, according to their *subject* (a *concept*), and

 defining a notion of preference for each module/concept based on the lexicographic closure semantics;

defining the semantics of the knowledge base by combining of the multiple preferences into a single preference.

Motivations

This multi-preferential approach provides a *spectrum of alternative semantics* depending on:

- the granularity of modularization (modules containing all conditionals vs. modules containing one conditional);
- the choice of the semantics for modules (here lexicographic closure semantics, but any ranked semantics could be adopted);
- how multiple preferences are combined into a single one.

This approach leads to a *preferential semantics* of the KB and a notion of *conditional entailment* which:

- satisfies the KLM properties of system P;
- is not subject to "blockage of property inheritance";
- deals properly with "*ambiguity preservation*" [Geffner and Pearl 92].

The Description Logic \mathcal{ALC}

Language of \mathcal{ALC}

Let N_C , N_R , N_I be the set of concept names, role names and individual names.

 \mathcal{ALC} concepts:

 $C := A \mid \top \mid \bot \mid \neg C \mid C \sqcap C \mid C \sqcup C \mid \forall S.C \mid \exists S.C$

where $A \in N_C$, $R \in N_R$

a first order logic or a polymodal logic (Schild,1991)

A Knowledge Base is a pair KB = (TBox, ABox):

- ► TBox contains a finite set of *inclusion axioms* C ⊆ D Mother_of_a_Doctor ⊆ ∃hasChild.Doctor.
- ► ABox is a set of *individual assertions* of the form C(a) and R(a, b), where $a, b \in N_l$, a set of individual names. For instance:

Female(mary), hasFriend(mary, carlo) (Italian □ ∃hasFriend.Engineer)(carlo)

ALC Semantics

An ALC interpretation is any structure $I = (\Delta^{I}, \cdot^{I})$ where:

- Δ' is a domain;
- ·¹ is an interpretation function that maps
 - each concept name A to set $A' \subseteq \Delta'$,
 - each role name *R* to a binary relation $R' \subseteq \Delta' \times \Delta'$,
 - each individual name *a* to an element $a' \in \Delta'$.
- ► .¹ is extended to complex concepts as follows:
 - $\succ \ \top' = \Delta \qquad \perp' = \emptyset \qquad (\neg C)' = \Delta C'$ $\succ \ (C \sqcap D)' = C' \land D'$

$$(C \sqcap D)' = C' \cap D' \qquad (C \sqcup D)' = C' \cup D$$

Satisfiability

An interpretation $\mathcal{M} = \langle \Delta, \cdot^{I} \rangle$ satisfies:

- a concept inclusion axiom $C \sqsubseteq D$ if $C' \subseteq D'$;
- an individual assertion C(a) if $a' \in C'$;
- ► an individual assertion R(a, b) if $(a', b') \in R_{a}^{l}$

Preferential extensions of DLs

- Preferential extensions of description logics allow defeasible inclusions in the knowledge base to model typical properties of individuals. Kraus Lehmann and Magidor's conditional assertions C ~ D become, for ALC:
 - typicality inclusions T(C) ⊑ D (Giordano et al., LPAR 2007, FI 2009) based on the preferential semantics [KLM 90];
 - ► defeasible inclusions C ≂ D (Britz et al. KR 2008) based on the rational semantics [LM 92].

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ALC with typicality

Preferential Interpretations

A preferential interpretation is a structure $\langle \Delta, <, \cdot' \rangle$ where:

- ► △ and ·¹ are a domain and an interpretation function, as in ALC interpretations;
- ► < is an irreflexive and transitive relation over △ and is well-founded.

Basic idea: x < y means: x is more normal than y

- $(\mathbf{T}(C))^{\prime} = Min_{<}(C^{\prime})$
- $\mathcal{M} \models \mathbf{T}(\mathcal{C}) \sqsubseteq \mathcal{D} \text{ iff } (\mathbf{T}(\mathcal{C}))^{\prime} \subseteq \mathcal{D}^{\prime}$

Ranked interpretations

modularity: for all $x, y, z \in \Delta$, if x < y then either x < z or z < yEach $x \in \Delta$ has a rank $k_{\mathcal{M}}(x)$, where $k_{\mathcal{M}} : \Delta \rightarrow Ord$

Entailment

► T(C) □ D is preferentially (rationally) entailed by K, if it is satisfied in all preferential (ranked) models M of K.

A minimal model semantics for RC

Preferential and rational entailment define a weak notion of inference.

Prefer the ranked models which minimize the rank of individuals: (Casini, et al., DL 2013), (Giordano, et al., DL 2013) *Minimal canonical ranked models* of the KB provide a semantic characterization *of the rational closure* of *ALC*.

It is a generalization of DLs of the canonical model semantics of Rational Closure by Lehmann and Magidor (1992).

The first Rational Closure construction for DLs was developed for ALC by Casini and Straccia (2010).

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Rational Closure ranking

RC construction assigns a rank to each defeasible inclusion and to each concept: less exceptional concepts have lower rank.

Example

 $PhDSt \sqsubseteq St$ $Emp \sqsubseteq Adult$ $PhDSt \sqsubseteq Adult$ $PrimarySchoolSt \sqsubseteq Children$



Advantages and drawbacks of the rational closure

- The rational closure has good computational properties. and can be extended to lightweight and expressive DLs
- Too weak: All or nothing: "the blocking of property inheritance problem" [Pearl,90]
 "the drowning problem" [Benferhat,Dubois,Prade,93].
- On the other hand: some conclusions are too strong [Geffner and Pearl,92]
- Exploits a unique preference relation < among individuals, but bob < Student tom and tom < Employee bob.</p>

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Alternative constructions to Rational Closure:

- ► The lexicographic closure, originally introduced by Lehmann and, for ALC, by (Casini and Straccia, 2012),
- The relevant closure (Casini, Meyer, Moodley, Nortje, 2014)
- ► The logic of overriding DL^N (Bonatti, Faella, Petrova, Sauro, 2015)
- Skeptical closure (Pruv, 18), MP-closure (ECSQARU'19)

Multipreferences in DLs

- (Bonatti Lutz, Wolter, 2009) circumscriptive KBs also allow abnormal instances of a class C with respect to a given aspect P using binary abnormality predicates Ab(P, x)
- ► (Fernandez Gil, 2014) several typicality operators T₁, T₂,... and preference relations <₁, <₂,... in ALC + T_{min}
- ► (Gliozzi, 2016) multiple (ranked) preferences associated to aspects (concepts) <_{A1}, <_{A2},...; a refinement of RC

Modular multi-concept knowledge bases

A modular multi-concept knowledge base *K* is a tuple $\langle T, D, m_1, \ldots, m_k, A, s \rangle$, where:

- T is an ALC TBox,
- ► D is a set of typicality inclusions, such that $m_1 \cup \ldots \cup m_k = D$,
- \mathcal{A} is an ABox, and
- ► *s* is a function associating each module m_i with a concept: if $s(m_i) = C_i$, then C_i is the subject of m_i .

Example

Let *K* be the knowledge base $\langle \mathcal{T}, \mathcal{D}, m_1, m_2, m_3, \mathcal{A}, s \rangle$, where $\mathcal{A} = \{St(mary), Emp(tom), PhDSt(bob)\},\$

TBox \mathcal{T}

 $PhDSt \sqsubseteq St$ $Emp \sqsubseteq Adult$ $PhDSt \sqsubseteq Adult$ $PrimarySchoolSt \sqsubseteq Children$

Module m_1 with subject *Student*; $s(m_1) = St$

 $\begin{array}{ll} \mathbf{T}(St) \sqsubseteq \neg Has_Scholarship \\ \mathbf{T}(PhDSt) \sqsubseteq Has_Scholarship \\ \mathbf{T}(Emp \sqcap St) \sqsubseteq Busy \\ \mathbf{T}(Emp \sqcap St \sqcap OnLeave) \sqsubseteq \neg Busy \end{array} \qquad \begin{array}{l} \mathbf{T}(St) \sqsubseteq Young \\ \mathbf{T}(PhDSt) \sqsubseteq Bright \\ \mathbf{T}(Emp \sqcap St) \sqsubseteq Busy \\ \mathbf{T}(St) \sqsubseteq St \sqcap St \sqcap St \\ \mathbf{T}(St) \sqsubseteq St \\ \mathbf{T}(St) \\ \mathbf{T}(St) \sqsubseteq St \\ \mathbf{T}(St) \sqsubseteq St \\ \mathbf{T}(St) \sqsubseteq \\ \mathbf{T}(St) \sqsubseteq \\ \mathbf{T}(St) \sqsubseteq \\ \mathbf{T}(St) \\ \mathbf{T}(St)$

Module m_2 with subject *Employee*; $s(m_2) = Emp$ $T(Emp) \sqsubseteq \neg Young$ $T(Emp \sqcap St) \sqsubseteq Busy$ $T(Emp \sqcap St \sqcap OnLeave) \sqsubseteq \neg Busy$

Multipreference semantics for modular KBs

We introduce a preferential relation $<_i$ for each module m_i .

We exploit the *lexicographic closure semantics* which already accounts for *specificity within the modules*.

Another option for small modules, one for each distinguished concept C_i , containing "only" the typicality inclusions $T(C_i) \sqsubseteq D$: user defined ranks in ranked \mathcal{EL}^+_{\perp} KBs [ICLP2020].

Module m_1 with subject *St* (Student), contains:

$\mathbf{T}(St) \sqsubseteq \neg Has_Scholarship$	rank = 0
$\mathbf{T}(St) \sqsubseteq Young$	rank = 1
$\mathbf{T}(St) \sqsubseteq \exists has_Classes.T$	rank = 2

Module *m*₂ with subject *PhDSt* (PhD Student), contains:

$\mathbf{T}(PhDSt) \sqsubseteq Has_Scholarship$	rank = 0
$\mathbf{T}(PhDSt) \sqsubseteq Bright$	rank = 1

In this case, treatment of *specificity among modules* is required.

Lexicographic multipreferences for modular KBs

The lexicographic closure (LC)

- Introduced by Lehmann (1995) in terms of maxiconsistent sets and in terms of preferential models;
- Casini and Straccia have defined the lexicographic closure for ALC (2012) based on maxiconsistent sets.

Lehmann's preferential semantics of LC

A preference relation < on the set of propositional interpretations (a *modular partial order relation*):

w < w' iff $V(w) \prec V(w')$

where V(w) is the set of defaults violated in w and \prec is a seriousness ordering among sets of defaults. (the defaults violated in w are less serious than the defaults violated in w'). The lexicographic closure semantics for \mathcal{ALC}

Let $\langle \Delta, \cdot^I \rangle$ be a (finite) \mathcal{ALC} model of *K*.

We define a preferential (ranked) model $\mathcal{N} = \langle \Delta, \langle \cdot, \cdot' \rangle$ for \mathcal{K} .

Preference relation < on Δ We define a *(modular) partial order relation* < on Δ as follows: x < y *iff* $V(x) \prec V(y)$ where V(x) is *the set of defaults violated in x* and \prec is a *seriousness ordering* among sets of defaults. (the defaults violated in *x* are less serious than the defaults violated in *y*).

Note: < as in Lehmann's preference semantics but preferences are over Δ .

Lexicographic models for module m_i

The projection of the knowledge base *K* on module m_i as the knowledge base $K_i = \langle T \cup m_i, A \rangle$.

- A lexicographic model of K_i
- is a *preferential model* $\mathcal{N}_i = \langle \Delta, \langle \cdot, \cdot^I \rangle$ of K_i , such that:
 - $\langle \Delta, \cdot' \rangle$ is an \mathcal{ALC} model of $\langle \mathcal{T}, \mathcal{A} \rangle$ and
 - satisfies the following condition:

$$x <_i y \quad \text{iff } V_i(x) \prec_i V_i(y). \tag{1}$$

where $V_i(x)$ is the set of *defaults in* K_i *violated by* x.

A lexicographic model $\mathcal{N}_i = \langle \Delta, <_i, \cdot^I \rangle$ of $K_i = \langle \mathcal{T} \cup m_i, \mathcal{A} \rangle$ is a *ranked model* of K_i .

Multiconcept lexicographic models

A multi-concept model for *K* is a *multi-preference interpretation* with *a preference relation* $<_i$ for each module m_i . Let $K = \langle T, D, m_1, ..., m_k, A, s \rangle$ be a multi-concept knowledge base.

Multi-concept lexicographic models

A multi-concept lexicographic model $\mathcal{M} = \langle \Delta, <_1, \dots, <_k, \cdot^l \rangle$ of K is a multi-concept interpretation for K, such that, for all $i = 1, \dots, k, \ \mathcal{N}_i = \langle \Delta, <_i, \cdot^l \rangle$ is a lexicographic model of $K_i = \langle \mathcal{T} \cup m_i, \mathcal{A} \rangle$.

We will consider *canonical* multi-concept lexicographic models of K, where the domain Δ contains, roughly speaking, *as many domain elements as consistent with* K.

Reasoning in a multi-concept knowledge base

In a multi-concept lexicographic model $\mathcal{M} = \langle \Delta, <_1, \dots, <_k, \cdot^I \rangle$ of *K*, preference relation $<_i$ can be used to answer *queries over module* m_i (i.e. queries with subject C_i).

The query "Are all typical Phd students young?" can be evaluated in module m_2 with subject Student by verifying that in all canonical multiconcept models of K,

 $min_{\leq_2}(PhD_Student^l) \subseteq Young^l$.

The answer would be positive, as the property of students of being normally young is inherited by PhD Student.

But what about the query:

"Are typical employed students young?" Employed students are concerned with both subjects *Student* and *Employee*, i.e., with both modules m_1 and m_2 , and with both perferences $<_1$ and $<_2$.

Combining preferences

We need to *combine the preference relations* $<_i$ into a single one <.

Let \leq_i be defined as follows: $x \leq_i y$ iff $y \not<_i x$. As $<_i$ is a modular partial order, \leq_i is a total preorder.

Given a canonical multi-concept lexicographic model $\mathcal{M} = \langle \Delta, <_1, \ldots, <_k, \cdot^l \rangle$ of *K*, we define a *global preference relation* < on Δ by Pareto combination:

$$x < y$$
 iff (i) for some $i = 1, ..., k$, $x <_i y$ and (*)
(ii) for all $j = 1, ..., k, x \leq_j y$,

< *is a partial order relation* (but, in general, modularity does not hold for <).

 $\mathcal{M}^{\mathbf{P}} = \langle \Delta, <, \cdot^{I} \rangle$ is a combined lexicographic model of K.

Combined lexicographic models of K and entailment

A combined lexicographic model $\mathcal{M}^{\mathsf{P}} = \langle \Delta, <, \cdot' \rangle$ of \mathcal{K} :

- is a preferential interpretation;
- ► satisfies strict inclusions and assertions in K;
- ▶ but is *not* required to satisfy *all typicality inclusions* $\mathbf{T}(C) \sqsubseteq D$ in K.

Example

 $\begin{array}{l} s(m_1) = Student \\ \textbf{T}(Student) \sqsubseteq YoungPerson \\ s(m_2) = YoungPerson: \\ \textbf{T}(YoungPerson) \sqsubseteq Student \\ \textbf{If Bob and John be young persons and also students, and \\ bob <_1 john and john <_2 bob (they are incomparable wrt <) \end{array}$

It may be that:

 $bob, john \in min_{<}(Student^{I}) \not\subseteq min_{< 1}(Student^{I})$ (note that this KB has no KLM style preferential model)

Multi-concept lexicographic entailment

To require that all typicality inclusions in *K* are satisfied in $\mathcal{M}^{\mathbf{P}}$, the notion of m_l^c -model of *K* can be strengthened by defining **T**-compliant m_l^c -models of *K*, i.e., m_l^c -models of *K* satisfying all typicality inclusions in *K*.

The notions of m_l^c -entailment and **T**-compliant m_l^c -entailment can be defined in the obvious way.

- satisfy the KLM postulates of preferential consequence relations, which can be reformulated for a typicality logic, considering that typicality inclusions T(C) ⊑ D stand for conditionals C ∼D.
- avoid the drowning problem (as lexicographic closure does)
- avoid unwanted inferences of RC and LC (if proper modularization)

Avoiding too strong conclusions of RC

- The refinements of RC (including MP-closure) avoid the "inheritance blocking" but allow for too strong conclusions (Geffner and Pearl, 92).
- To deal correctly with ambiguity preservation Conditional Entailment (Geffner and Pearl, 92) and system LCD (Benferhat, Saffiotti and Smets, 2000) abandon Rational Monotonicity (RM). System JLZ (Weydert, 2003) exploits a canonical ranking construction and verifies (RM).
- Combining multiple preferences is a simple approach which may avoid too strong conclusions (depending on the definition of specificity ≺) as well as inheritance blocking. This is the case for the concept-wise multipreference semantics.

Example: ambiguity preservation

A reformulation of an example by Geffner and Pearl (AIJ'92).

Consider module m_1 with subject *Citizen*:

 $\mathbf{T}(Italian) \sqsubseteq DriveFast$

 $\mathbf{T}(\mathit{Italian}) \sqsubseteq \mathit{HomeOwner}$

and module *m*₂ with subject *Student*:

 $\mathbf{T}(PhDStudent) \sqsubseteq \neg HomeOwner$

 $T(PhDStudent) \sqsubseteq Has_Scholarship$

In RC and LC we would neither conclude that Italian PhD students are home owners nor that they are not home owners. If the inclusion

 $T(Student) \sqsubseteq \neg Has_Scholarship$

is added, being *PhDStudent* \sqsubseteq *Student*, RC and LC conclude:

 $T(Italian \sqcap PhDStudent) \sqsubseteq \neg HomeOwner$

("the conflict is resolved anomalously").

Separating typicality inclusions in the two modules m_1 and m_2 avoids this conclusion.

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A hierarchy of modules

We may have introduce different modules

 m_1, m_2, m_3, m_4

with subject (resp.)

Student, Employee, PhDStudent, Student \sqcap Employee. A specificity relation \succ among concepts (modules) has to be taken into account:

PhDStudent > Student

(PhDStudent is more specific than Student)

Student

Employee

Student

Student
¬ Employee > Employee

Specificity is needed for defining the global preference <. We can adopt the *modified Pareto condition* from [TPLP2020].

$$x < y$$
 iff (i) $x <_i y$, for some $C_i \in C$, and
(ii) for all $C_j \in C$, $x \leq_j y$ or $\exists C_h(C_h \succ C_j \text{ and } x <_h y)$

The preference relation $<_3$ for PhD students overrides the preference relation $<_1$ for Students.

Conclusions

- Combining multiple preferences is a simple approach which may avoid too strong conclusions of RC as well as inheritance blocking. This is the case for the concept-wise multipreference semantics.
- The approach to preference combination has different instantiations depending on: the granularity of aspects (e.g., concepts or modules), the definition of preferences <_{Ai} and the notion specificity [NMR2020].
- An algebraic framework for preference combination has been proposed by Bozzato et al. [TPLP2021].
- Plausibility of a concept-wise semantics has been supported [Cilc 2020] by showing that it can be used to define a logical interpretation of self-organising maps (SOMs), psychologically and biologically plausible neural network models (Kohonen'01).

Related work

- Brewka's Preferred subtheories (1998) and framework of Basic Preference Descriptions (2004)
- Preference fusion to define system ARS (Kern-Isberner and Ritterskamp, 2010) a refinement System Z;
- Syntax Splitting of conditional KBs for c-representations (Kern-Isberner, Beierle, Brewka, 2020);
- ► Gil (2014) has defined a multipreference formulation of the typicality DL ALC + T_{min} to avoid inheritance blocking;
- Gliozzi's multipreference semantics (2016) to define a refinement of rational closure.
- Britz and Varzinczak (2018, 2019) associate preference relations to roles, to define defeasible role quantifiers and defeasible role inclusions;
- Delgrande and Rantsoudis (NMR 2020) have proposed a multi-preferential approach for representing defaults in FOL

Related work

- In the CKR (Contextualized Knowledge Repositories) framework, by Bozzato, Eiter and Serafini (2014,2018), defeasible axioms are allowed and exceptions can be handled by overriding. Extended to general contextual hierarchies (2019). ASP based reasoning for SROIQ-RL.
- In [ICLP 2020] the approach is applied to ranked *EL*⁺_⊥ knowledge bases, where each module corresponds to a concept *C_i* and only contains defaults **T**(*C_i*) ⊑ *D* (with their ranks). ASP (and asprin) used for defeasible reasoning in ranked *EL*⁺_⊥ KBs.
- An algebraic framework for preference combination has been proposed by Bozzato, Eiter and Kiesel [TPLP2021].
- Weighted defeasible DL knowledge bases for modeling Multilayer Perceptrons under a fuzzy multipreference semantics [JELIA 2021]. For the two-valued preferential semantics ASP based reasoning in ICLP2021.

Thank you!!!!!

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